

# Problem sheet 2

Due Monday, April 30th at 17:30.

## Question 1

Consider a graph  $G$  on  $n$  vertices and  $m$  edges. Prove that  $G$  contains at least  $\frac{4m}{3n}(m - \frac{n^2}{4})$  triangles.

**Solution** Let  $G$  and  $|V(G)| = n, |E(G)| = m$ . Let  $x, y \in V(G)$ , then  $|N(x) \cap N(y)| \geq d(x) + d(y) - n$ . Therefore, in particular, any edge  $\{x, y\}$  is contained in at least  $d(x) + d(y) - n$  triangles. Hence the total number of triangles in  $G$  is at least

$$\frac{1}{3} \sum_{\{x,y\} \in E(G)} (d(x) + d(y) - n) = \frac{1}{3} \sum_{x \in V(G)} d^2(x) - \frac{1}{3}nm \geq \frac{\left(\sum_{x \in V(G)} d(x)\right)^2}{3n} - \frac{1}{3}nm = \frac{4m^2}{3n} - \frac{1}{3}nm.$$

## Question 2

Let  $G$  be a graph on  $n$  vertices and let  $N_t, t \geq 2$ , denote the number of copies of  $K_t$  in  $G$ . Prove that

$$\frac{N_{t+1}}{N_t} \geq \frac{1}{t^2 - 1} \left( t^2 \frac{N_t}{N_{t-1}} - n \right).$$

(**Hint:** Double count pairs  $(A, U)$  such that  $|A| = |U| = t, |A \cap U| = t - 1$ ,  $A$  induces  $K_t$  and  $U$  induces a non-complete graph of  $t$  vertices.)

**Solution** Let  $A_1, A_2, \dots, A_{N_t}$  be the collection of different  $K_t$  in  $G$  and let  $B_1, B_2, \dots, B_{N_{t-1}}$  be the collection of different  $K_{t-1}$  in  $G$ . Let  $a_i$  be the number of copies of  $K_{t+1}$  in  $G$  containing  $A_i$ . Let  $b_i$  be the number of copies of  $K_t$  in  $G$  containing  $B_i$ .

We double count the number  $N$  of pairs  $(A, U)$  where  $A \subseteq V(G)$  induces  $K_t$  and  $U \subseteq V(G)$  induces a non-complete graph on  $t$  vertices and  $|A \cap U| = t - 1$ . Considering, for the  $i^{\text{th}}$  copy  $A_i$  of  $K_t$ , pairs  $(A_i, U_j(x))$ , where  $U_j(x) \subset A_i \cup \{x\}$  is a set containing a vertex  $x \notin A_i$  and a vertex  $y \in A_i$  that is not adjacent to  $x$ ,  $1 \leq j \leq t - 1$ , yields

$$N \geq \sum_{i=1}^{N_t} (t-1)(n-t-a_i).$$

This is true because there are at least  $(n - t - a_i)$  vertices in  $G$  which have a non-neighbour in  $A_i$ . After choosing such a vertex  $v$  with a non-neighbour  $w$  in  $A_i$ , there are at least  $(t - 1)$  to remove one of the vertices in  $A_i$  which is not  $w$ .

Also we have,

$$N = \sum_{i=1}^{N_{t-1}} b_i(n - (t - 1) - b_i).$$

We have  $\sum_{i=1}^{N_t} a_i = (t + 1)N_{t+1}$  and  $\sum_{i=1}^{N_{t-1}} b_i = tN_t$ . Using Jensen's inequality (for the concave function  $f(x) = x(n - t + 1 - x)$ ) we show

$$(t-1)(n-t)N_t - (t^2-1)N_{t+1} \leq N \leq N_{t-1} \frac{\sum_{i=1}^{N_{t-1}} b_i}{N_{t-1}} \left( n - t + 1 - \frac{\sum_{i=1}^{N_{t-1}} b_i}{N_{t-1}} \right) = t(n-t+1)N_t - \frac{t^2 N_t^2}{N_{t-1}}.$$

The above gives us the required inequality.

### Question 3

Let  $H$  be a graph with  $\chi(H) = t$ .

- (a) Prove that  $\text{ex}(n, H) \geq \text{ex}(n, K_t)$  for all  $n \geq 1$ .
- (b) Suppose that  $\text{ex}(n, H) = \text{ex}(n, K_t)$  for some  $n \geq t$ .

Prove that there is an edge  $e$  in  $H$  such that  $\chi(H - e) < t$ .

### Solution

- (a) The complete multi-partite graph  $G$  with parts of sizes as equal as possible has chromatic number  $t-1$ , hence can not contain  $H$  as a subgraph. This implies that  $\text{ex}(n, H) \geq |E(G)|$ . By Turán's theorem  $|E(G)| = \text{ex}(n, K_t)$ .
- (b) Let  $G$  be a graph as in part (a). Since  $n \geq t$ , the graph  $G$  is not a complete graph. Let  $G'$  be a graph obtained from  $G$  by adding an edge  $e'$  within one of the parts. Then  $G'$  contains a copy  $K$  of  $H$ , since  $\text{ex}(n, H) = \text{ex}(n, K_t) < |E(G')|$ . Since  $G$  does not contain  $H$ , the edge  $e'$  is contained in  $K$ . Let  $e$  denote the edge in  $H$  corresponding to  $e'$ . Then  $H - e$  is contained in  $G$  and therefore  $\chi(H - e) \leq \chi(G) = t - 1 < t$ .

### Question 4

Suppose that  $H$  is a graph with  $\text{ex}(n, H) \leq \lambda \binom{n}{2}$  for some constant  $\lambda$ ,  $0 < \lambda < 1$ , and  $n \geq n_0$ . Prove that for any  $\epsilon > 0$  and sufficiently large  $n$  any graph on  $n$  vertices and  $(\lambda + \epsilon) \binom{n}{2}$  edges contains at least  $c(\epsilon, n_0)n^{|V(H)|}$  copies of  $H$ .

**Solution** Let  $G$  be a graph on  $n$  vertices and at least  $(\lambda + \epsilon) \binom{n}{2}$  edges. Let  $t$  be the number of sets  $N$  on  $n_0$  vertices and more than  $\lambda \binom{n_0}{2}$  edges in  $G[N]$ .

We double-count pairs  $(e, N)$ , where  $N$  is a set of  $n_0$  vertices in  $G$  and  $e$  is an edge with both endpoints in  $N$ , let  $P$  be the number of such pairs. Then by first considering the edges and then choosing additional  $n_0 - 2$  vertices to  $N$  we get that,

$$P = |E(G)| \binom{n-2}{n_0-2} \geq (\lambda + \epsilon) \binom{n}{2} \binom{n-2}{n_0-2} = (\lambda + \epsilon) \binom{n}{n_0} \binom{n_0}{2}.$$

On the other hand, by first considering sets of  $n_0$  edges and then the edges in those sets we get,

$$P \leq t \binom{n_0}{2} + \left( \binom{n}{n_0} - t \right) \lambda \binom{n_0}{2} = t(1 - \lambda) \binom{n_0}{2} + \lambda \binom{n}{n_0} \binom{n_0}{2}.$$

This is true because in  $t$  sets  $N$  on  $n_0$  vertices we can have at most  $\binom{n_0}{2}$  edges. On the rest  $\left( \binom{n}{n_0} - t \right)$  we have at most  $\lambda \binom{n_0}{2}$  edges.

These inequalities yield

$$t \geq \frac{\epsilon}{(1 - \lambda)} \binom{n}{n_0} \geq \epsilon \binom{n}{n_0}.$$

For each set  $N$  on  $n_0$  vertices and more than  $\lambda \binom{n_0}{2}$  edges there is a copy of  $H$  in  $G[N]$ . The vertices of each copy of  $H$  in  $G$  are contained in  $\binom{n - |V(H)|}{n_0 - |V(H)|}$  sets on  $n_0$  vertices. Let  $v = |V(H)|$ . Therefore the number of copies of  $H$  in  $G$  is at least

$$t \binom{n-v}{n_0-v}^{-1} \geq \epsilon \binom{n}{n_0} \binom{n-v}{n_0-v}^{-1} = \epsilon \frac{n!(n_0-v)!}{n_0!(n-v)!} \geq \epsilon \frac{(n_0-v)!}{n_0!} c(v) n^v.$$

For the last inequality we used that there is a constant  $c(v)$ , depending on  $v$ , such that  $\frac{n!}{(n-v)!} \geq c(v) n^v$ . This is given by Stirling's approximation, i.e., there are constants  $c_1$  and  $c_2$  such that for all  $n > 0$  we have  $c_1 \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \leq n! \leq c_2 \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$ . This yields

$$\frac{n!}{(n-v)!} \geq \frac{c_1 \sqrt{2\pi n} \left(\frac{n}{e}\right)^n}{c_2 \sqrt{2\pi(n-v)} \left(\frac{n-v}{e}\right)^{n-v}} \geq \frac{c_1}{c_2} e^v n^v.$$